A Survey on Algorithm Design Techniques

A.Mahalakshmi¹, Devi selvam², L.Viji³

^{1, 2} Department of CSE ³ Department of IT ^{1, 2} Sri Shakthi Institute of Engineering and Technology,Tamilnadu,India ³ Velalar College of Engineering and Technology,Tamilnadu,India

Abstract- Algorithms are recognized as steppingstone of computing. This survey describes different algorithm design techniques such as Brute force, Greedy, Divide and Conquer, Dynamic programming, Backtracking, Branch and Bound and many more. It describes how a particular technique is used for a specific problem by considering various parameters. This survey compares different algorithm design strategies along with applications.

Keywords- Brute force, Greedy, Divide and Conquer, Dynamic programming, Backtracking, Branch and Bound.

I. INTRODUCTION

Algorithm is a step-by-step procedure, which defines a set of instructions to be executed in a certain order to get the desired output for any valid input. There exist a number of algorithms, every algorithm is problem specific. The choice of an algorithm may not just depend on computational complexity; it also depends upon the characteristics, advantages and disadvantages. This report shows how an algorithm is best for a particular situation, based upon their advantages and comparison with others. Some algorithm design strategies are problem specific means they are well suited for some specific problem and have disadvantage against another problem. The basic question here is: How to choose the approach? First, by understanding the problem, and second, by knowing various problems and how they are solved using different approaches.



Figure 1.



II. PURPOSE OF ALGORITHM ANALYSIS

There are no well-defined standards for writing algorithms. Rather, it is problem and resource dependent. Algorithms are never written to support a particular programming code.We should know the problem domain, for which we are designing a solution.

The problem domain may belongs to one of the following categories

- 1. Sorting
- 2. Searching
- 3. Combinatorial Problems
- 4. Numerical Problems
- 5. Graph Problems

A. problem can be solved in more than one ways.



Figure 3.

Thus, we will have many solutions for a single issue. Among that to choose the best solution, we are in need of Analysis.

ALGORITHM COMPLEXITY



Figure 4.

Suppose A is an algorithm and n is the size of input data, the time and space used by the algorithm A are the two main factors, which decide the efficiency of A.

- Time Factor Time is measured by counting the number of basic operations, such as comparisons in the searching algorithm.
- Space Factor Space is measured by counting the maximum memory space required by the algorithm.
- Asymptotic analysis refers to computing the running time of any operation in mathematical units of computation. Usually, the time required by an algorithm falls under three types –
- Best Case Minimum amount of time required for program execution.
- Average Case Average time required for program execution.
- Worst Case Maximum amount of time required for program execution.

III. ALGORITHM DESIGN TECHNIQUES

Algorithm design technique is the systematic approach for solving issues under various domains. They provide templates suited to solving a broad range of diverse problems. They can be translated into common control and data structures provided by most high-level languages.

DESCRIPTION:

A BRUTE FORCE

Brute force is a problem solving approach in which we will "generate all possible solutions and select the optimal one ".It is a straightforward approach proceeds in simple and obvious way usually based directly on problem's statement and definitions of the concepts involved.

A. Exhaustive search

It is a brute-force approach, applied to solve combinatorial problems. The algorithm that tries every possible solution is known as exhaustive search. Also known as the British Museum algorithm. It can be used to reduce the search space.

B. GREEDY TECHNIQUE

A greedy algorithm, as the name suggests, always makes the choice that seems to be the best at that moment. A greedy algorithm always makes the next available best choice. Here at each step, decision made must be

- Feasible The solution should satisfy the given constraints
- Locally Optimal Best among all possible feasible solutions
- Irrevocable Once made, it cannot be changed

Example-Change Making Problem

This problem is to count to a desired value by choosing the least possible coins and the greedy approach forces the algorithm to pick the largest possible coin. If we are provided coins of Rs 1, 5, 10 and 20 and we are asked Rs 48 then the greedy procedure will be -

- 1 Select 3 coins of 1 Rs, the remaining count is 45
- 5 Then select one 5 Rs coin, the remaining count is 40
- 20 And finally, selection of two 20 Rs, solves the problem



Figure 5. Various algorithm design techniques are

C. DIVIDE AND CONQUER

There are three steps.

- Divide: Breaking the problem into several sub-problems that are similar to the original problem but smaller in size.
- Conquer: Solve the sub-problem recursively (successively and independently). If the sub problems are small enough, solve them in brute force fashion.
- Combine: Combine these solutions to sub-problems to create a solution to the original problem.

Divide and conquer algorithms have several appealing properties that make them a good match for modern parallel machines.



D. DYNAMIC PROGRAMMING

Dynamic programming is the method of solving the issue by overlapping of sub-problems.

Dynamic programming is used where we have problems, which can be divided into similar sub-problems, so that their results can be re-used. It uses the concept called "Principle of Optimality", which can be stated as" in a optimal sequences of choices, each subsequence must also be optimal".

There are two ways of doing this.

Page | 73

- 1.) Top-Down : Start solving the given problem by breaking it down. If you see that the problem has been solved already, then just return the saved answer. If it has not been solved, solve it and save the answer. This is referred to as Memorization.
- 2.) Bottom-Up : Analyze the problem and see the order in which the sub-problems are solved and start solving from the trivial subproblem, up towards the given problem. In

this process, it is guaranteed that the subproblems are solved before solving the problem. This is referred to as Dynamic Programming.

Knapsack problem

Given a set of n items, each with a weight and a value, determine the number of each item to include in a collection so that the total weight is less than or equal to a given limit and the total value is as large as possible.



E. BACKTRACKING

The principle idea of backtracking is to construct solutions on component at a time and evaluate such partially constructed candidates as follows. If a partially constructed solution can be developed further without violating the problem's constraints, it is done by taking the first remaining legitimate option for the next component. If there is no legitimate option for the next component, no alternatives for any remaining component need to be considered. In this case the algorithm backtracks to replace the last component of the partially constructed solution with its next option. Recur¬sion is the key concept in back¬track¬ing.

State space tree

This method is based on construction of state space tree. It is the tree organization of entire solution space(which has complete solution path and dead ends). Where the root represents the initial state before the search for a solution begins. The nodes at each nth level represent the choices for the nth component.

N Queens Problem

The N Queen is the problem of placing N chess queens on an $N \times N$ chessboard so that no two queens attack each other. For example, following is a solution for 4 Queen Problem.

State Space Tree





Figure 9.

F. BRANCH AND BOUND

In this method, based on the bound (Lower/Upper) values, the state space tree will be expanded.



We know that maximum valuable subset should be identified to solve knapsack problem. Hence, it comes under

upper bound. But in case of travelling salesman problem, minimum distance should be identified. Hence it belongs to lower bound.



Figure 12.

S salesman wants to visit all cities A, B, C and D. What is the best way to do this (cheapest airline tickets, and minimal travel time)?

Table 1.		
Algorithm Design Techniques	Applications	
Brute Force	Selection sort,	
	Bubble sort, String	
	matching and	
	Convex-hull	
	problems.	
Exhaustive	Travelling Salesman	
Search	Problem, Job	
	Assignment	
	Problem.	
Greedy	Graph Coloring,	
Technique	Huffman Coding,	
	Dijkstra's Shortest	
	Path	
	Algorithm, Container	
	Loading	
	Problem, Prim's &	
	Kruskal's Minimum	
	Spanning Tree	
	Algorithm.	
Divide and	Merge Sort, Quick	

Conquer	Sort, Binary Search,		
	Strassen's Matrix		
	Multiplication,		
	Closest pair and		
	Convex Hull		
	Problems.		
I. DYNAMIC	Computing		
PROGRAMMING	Binomial		
	Coefficient,		
	Knapsack Problem,		
	Floyd, Warshall,		
	Optimal Binary		
	Search Tree.		
Backtracking	N Queens Problem,		
	Subset Sum		
	Problem,		
	Hamiltonian Circuit		
	Problem		
Branch and	Knapsack Problem,		
Bound	Travelling Salesman		
	Problem, Job		
	Assignment		
	Problem		

Table 2.

Algorithm	Advantages	Disadvantages
Design		
Technique		
S		
	- High speed	-Very
		hardware
	- Very Simple	intensive.
Bruto	- Fairly simplistic	-Requires a lot
Force	attack	of processing
1 01 00	that doesn't	power.
	require a	
	lot of work to	
	setup or	
initiate.		
	-It's widely	-Cost of
	applicable.	generating
		candidate
Exhaustive	-It's simple.	solutions.
Search		
Startin	-Reliable and	-Impractical
	universal	(intractable) on
	method.	problems that
		are too big.

	-Easy to	-Don't work for
	implement.	some problems.
	F	F
	-Repeatedly taking	-Much slower.
Greedy	the next available	
Technique	best choice is	-It is not an
1	usually linear work.	automatic.
	-Much cheaper than	
	most other	
algorithms.		
	-Parallelism.	-Recursion
		Overhead.
	-Cache	
	Performance.	-For efficiency
		reasons,
Divide end	-Memory Access.	programs often
Conquer		use pointers
Conquer	-Used to solve	into arrays and
	difficult	pointer
	problems.	arithmetic to
		identify sub
	-High Efficiency.	problems.
		-Used only for
	111. -USED TO	overlapping of
		1 11
Π Οννά	AVOID	sub problems.
II. Dyna mic	AVOID CALCULATING	sub problems.
II. Dyna mic Programm	AVOID CALCULATING SAME STUFF TWICE	sub problems. -It's very
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula.
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so
II. DYNA MIC PROGRAMM ING	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive.
II. Dyna mic Programm ing	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive.
II. DYNA MIC PROGRAMM ING	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient
II. DYNA MIC PROGRAMM ING Backtracki	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE. -It is a step- by-step representatio n of a solution- very easy to understand. -It is	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from one state.
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from one state.
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from one state. -Requires large
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from one state. -Requires large amount of
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from one state. -Requires large amount of storage space
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from one state. -Requires large amount of storage space
II. DYNA MIC PROGRAMM ING Backtracki ng	AVOID CALCULATING SAME STUFF TWICE.	sub problems. -It's very tough task to identify recursive formula. -Multiple function calls are used, so expensive. -Inefficient when there is lots of branching from one state. -Requires large amount of storage space

	- Easy to	
	debug .	
	It is used to	Finding
Branch	-It is used to solve optimization problem. -It may traverse the tree in any manner, DFS or BFS.	-Finding pruning strategies require clever thinking technologies
	-It completely searches the state space tree to get more number of optimal solutions.	

III. METHODOLOGY

The main key of the study to make an attempt of solving the material handling station with the implementation of the discrete even simulation i.e. modeling the entire procedural steps exactly by using queues and the probability and statistic functions the simple model of 3 station material handling system is considered and the steps which as the intermediate process elements such as the machining centers inspection stations, packing, conveyor belts etc. are part of this study.

IV. COMPARISON BETWEEN A LGORITHM DESIGN TECHNIQUES

Table 3	i.
GREEDY TECHNIQUE	BRUTE FORCE
Very Fast	Slower than Greedy method.
Tries to find a localized optimum solution, which may eventually lead to globally optimized solutions	Used to solve small class of problems
BACKTRACKING	BRUTE FORCE

ISSN [ONLINE]: 2	2395-1052
------------------	-----------

In each step, this method	This method generates
checks- if it satisfies all the	all possible
conditions. If it does :	combinations and then
continues generating	checks if any of them is
subsequent solutions. If not :	the exact solution
goes one step backward to	
check for another path	

Table 4.

DIVIDE AND	DYNAMIC
CONQUER	PROGRAMMING
Inputs are divided at prespecified deterministic points.	It tries every possible split points.
Principle of Optimality concept is not used.	Principle of Optimality concept is used.
Solutions are combined to achieve an overall solution	Use the output of a smaller sub-problem and then try to optimize a bigger sub- problem.
 Memory access Divide-and-conquer algorithms naturally tend to make efficient use of memory caches. 	Here, each sub-problem is solved only once. There is no recursion. The key in dynamic programming is remembering.
2) The sub-problems are independent of each other.	Here, the sub-problems are not independent of each other.
D&C algorithms that are time-efficient often have relatively small recursion depth. <i>3</i>)	DP gives a better performance for many NP complete problems like TSP. Though the space needed is large, it reduces the complexity well.

Table 5.

BACKTRACKING		BRANCH AND BOUND	
A 12 11			
Applicable	to	both	Applicable only to
Optimization	and	Non-	Optimization problems.
optimization problems.			
State space tree is generated by		State space tree is	
using depth first search method.		generated by using	
			breadth first search
			method.
It don't consi	der any	bound	State space tree is
values.			generated based on
			bound values.

V. CONCLUSION

The main goal of this survey is to compare various algorithm design techniques based on the advantages, disadvantages and applications. Usually a given problem can be solved using various approaches however it is not wise to settle for the first that comes to mind. More often than not, some approaches result in much more efficient solutions than others. In future, the best algorithm design technique suitable for all types of problems can be identified.

REFERENCES

- Shailendra Nigam, Dr. Deepak Garg "Choosing Best Algorithm Design Strategies For a Particular Problem", In Proceedings of the IEEE International Advance Computing Conference (IACC 09), Thapar University Patiala, India (6-7 March 2009)
- [2] Exhaustive Search, Chapter 38 in Algorithms by Robert Sedgewick; Addison- Wesley, 1983.
- [3] A New Road Map of Algorithm Design Techniques by Anany Levitin in Dr. Dobb's Journal, April 2008.
- [4] Greedy Algorithms, Chapter 17 in Introduction to Algorithms by Thomas Cormen, Charles Leiserson, and Ronald Rivest, MIT Press, 1999.
- [5] Greedy Algorithms, Chapter 5 in Algorithms by Sanjoy Dasgupta, Christos Papadimitriou, and Umesh Vizirani, McGraw-Hill, 2008
- [6] http://ignou.ac.in/userfiles/SandeepFINAL_Unit1_Intro_2 1-03-2013.pdf
- [7] http://faculty.simpson.edu/lydia.sinapova/www/cmsc250/ LN250_Weiss/L28- Design.htm
- [8] https://gradeup.co/algorithm-design-techniques-i-6c2dec3e-c0da-11e5-abd0-9e294d44e4af
- [9] http://gdeepak.com/thesisme/Thesis-Choosing% 20Best% 20Algorithm% 20Design% 20Strategi es% 20For% 20a% 20Particular% 20Problem.pdf
- [10] https://en.wikipedia.org/wiki/Divide_and_conquer_algorit hm
- [11] http://stackoverflow.com/questions/30025421/differencebetween-backtracking-and-branch-and-bound

- [12] http://www.algorithmist.com/index.php/Exhaustive_Searc h
- [13] https://people.csail.mit.edu/rinard/divide_and_conquer/
- [14] [https://en.wikipedia.org/wiki/Bellman_equation